

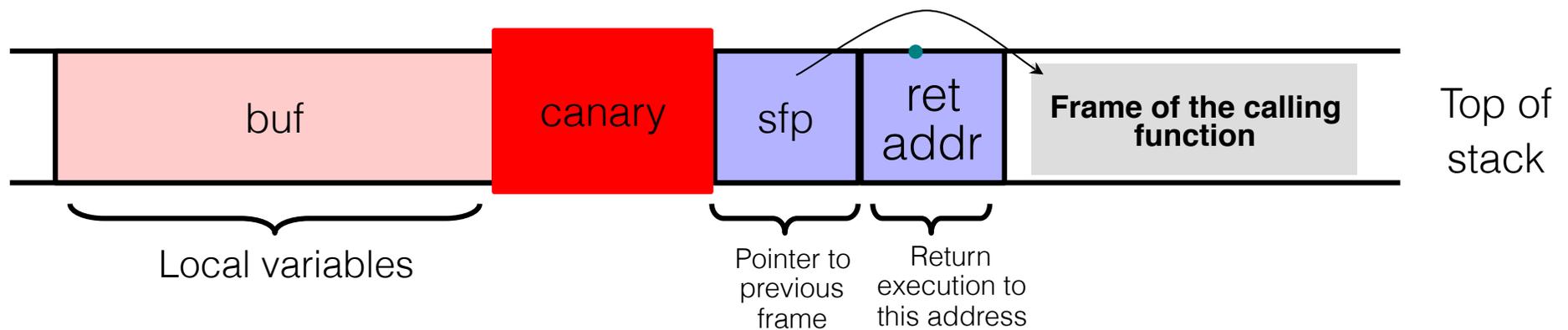
Integer Overflow, Format Strings

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Credits: Vitaly Shmatikov

Run-Time Checking: StackGuard

- Embed “canaries” (stack cookies) in stack frames and verify their integrity prior to function return
 - Any overflow of local variables will damage the canary



- Candidate Canaries
 - Choose random canary value picked on program start
 - Terminators: “\0”, newline, linefeed, EOF

What Can Still Be Overwritten?

- Other string buffers in the vulnerable function
- Any data stored on the stack

- Exception handling records

- Pointers to virtual method tables

C++: call to a member function passes as an argument “this” pointer to an object on the stack

Stack overflow can overwrite this object’s vtable pointer and make it point into an attacker-controlled area

When a virtual function is called (*how?*), control is transferred to attack code (*why?*)

Do canaries help in this case?

(Hint: when is the integrity of the canary checked?)

Example of Failed StackGuard — Litchfield's Attack

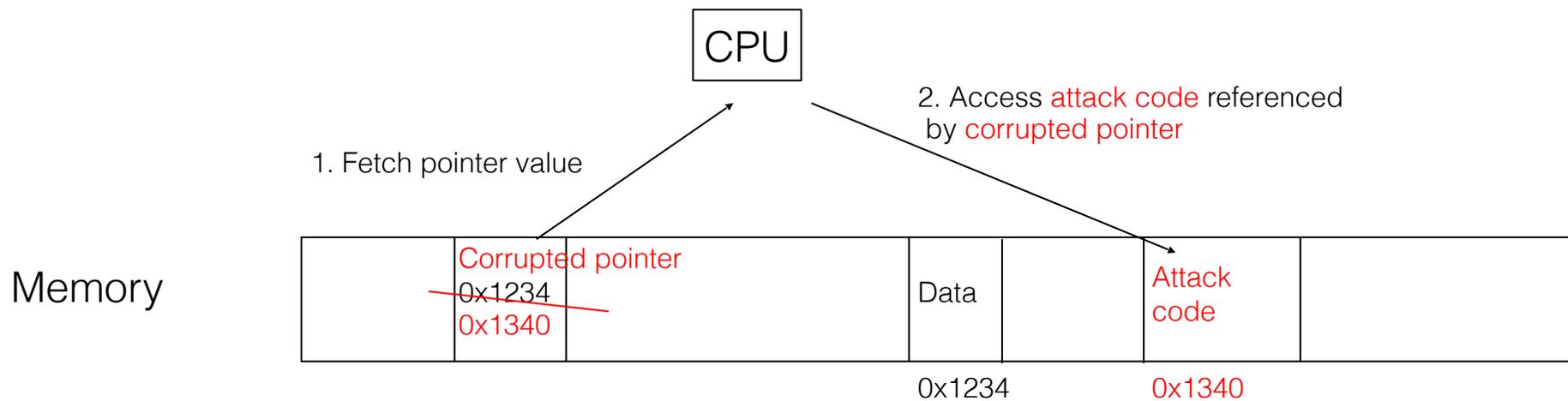
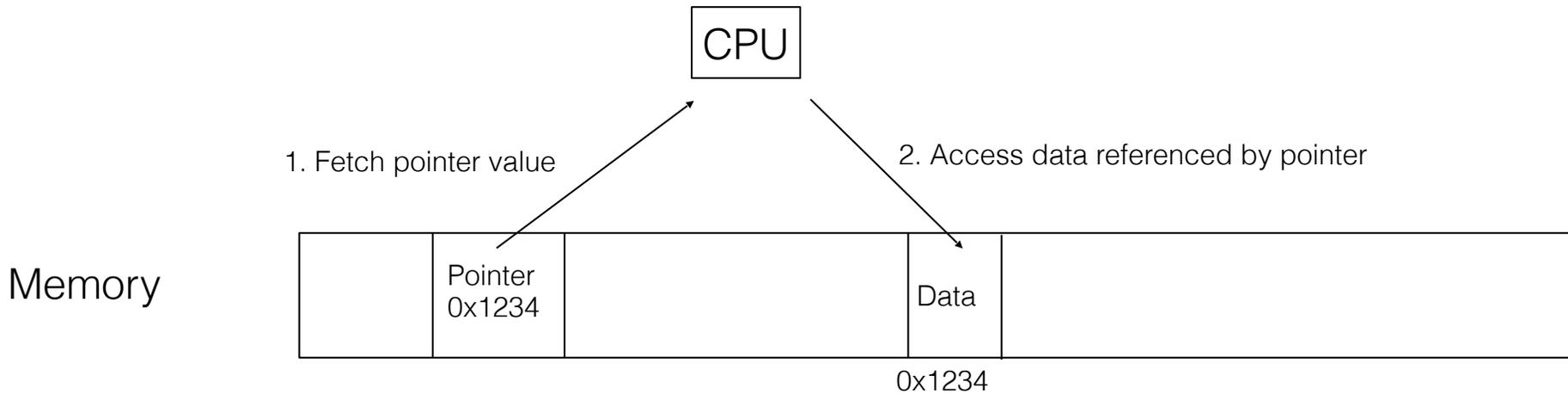
- Microsoft Windows 2003 server implements several defenses against stack overflow
 - Random canary (with /GS option in the .NET compiler)
 - When canary is damaged, exception handler is called
 - Address of exception handler stored on stack above RET
- Attack: smash the canary AND overwrite the pointer to the exception handler with the address of the attack code
 - Attack code must be on heap and outside the module, or else Windows won't execute the fake "handler"
 - Similar exploit used by CodeRed worm

PointGuard

- Attack: overflow a function pointer so that it points to attack code
- Idea: **encrypt all pointers** while in memory
 - Generate a random key when program is executed
 - Each pointer is XORed with this key when loaded from memory to registers or stored back into memory
 - Pointers cannot be overflowed while in registers
- Attacker cannot predict the target program's key
 - Even if pointer is overwritten, after XORing with key it will dereference to a "random" memory address

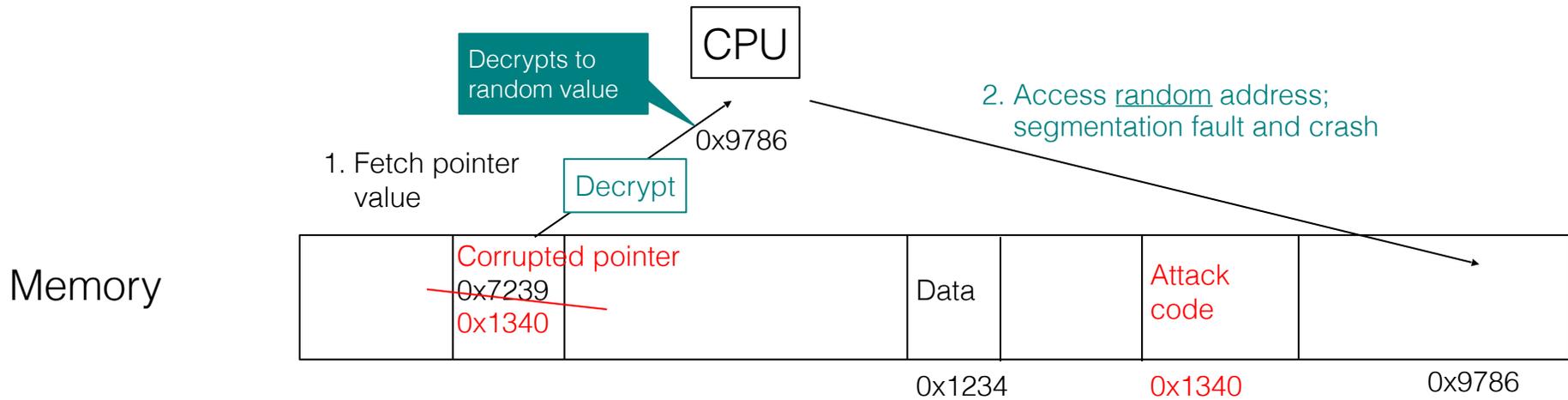
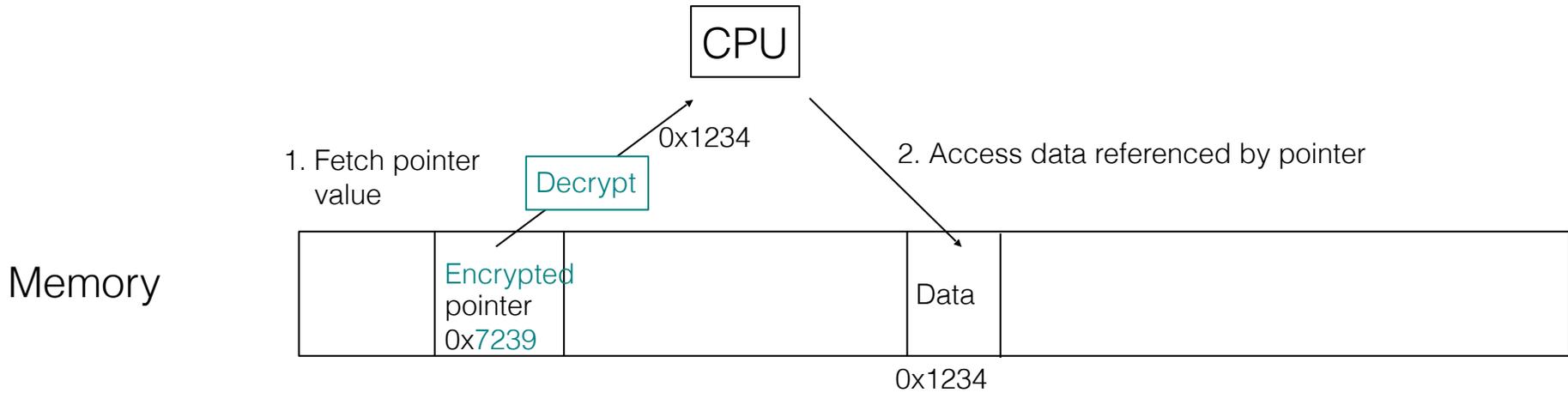
Normal Pointer Dereference

[Cowan]



PointGuard Dereference

[Cowan]

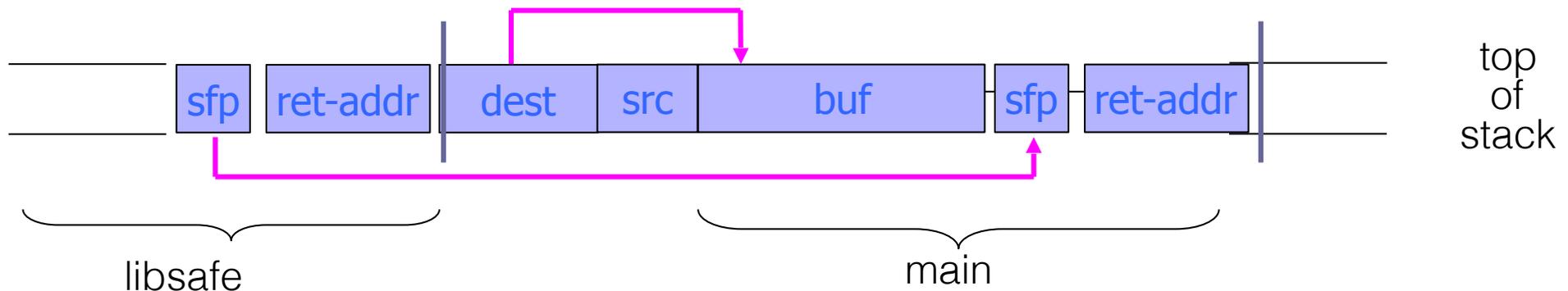


PointGuard Issues

- Must be very fast
 - Pointer dereferences are very common
- Compiler issues
 - Must encrypt and decrypt only pointers
 - If compiler “spills” registers, unencrypted pointer values end up in memory and can be overwritten there
- Attacker should not be able to modify the key
 - Store the key in a memory page inaccessible to adversaries
- PG'd code doesn't mix well with normal code
 - What if PG'd code needs to pass a pointer to OS kernel?

Libsafe

- Intercepts calls to `strcpy(dest, src)` and other unsafe C library functions
 - Checks if there is sufficient space in current stack frame $|\text{framePointer} - \text{dest}| > \text{strlen}(\text{src})$
 - If yes, does `strcpy`; else terminates application
- Dynamically loaded library – no need to recompile!



Limitations of Libsafe

- Protects frame pointer and return address from being overwritten by a stack overflow
- Does not prevent sensitive local variables below the buffer from being overwritten
- Does not prevent overflows on global and dynamically allocated buffers

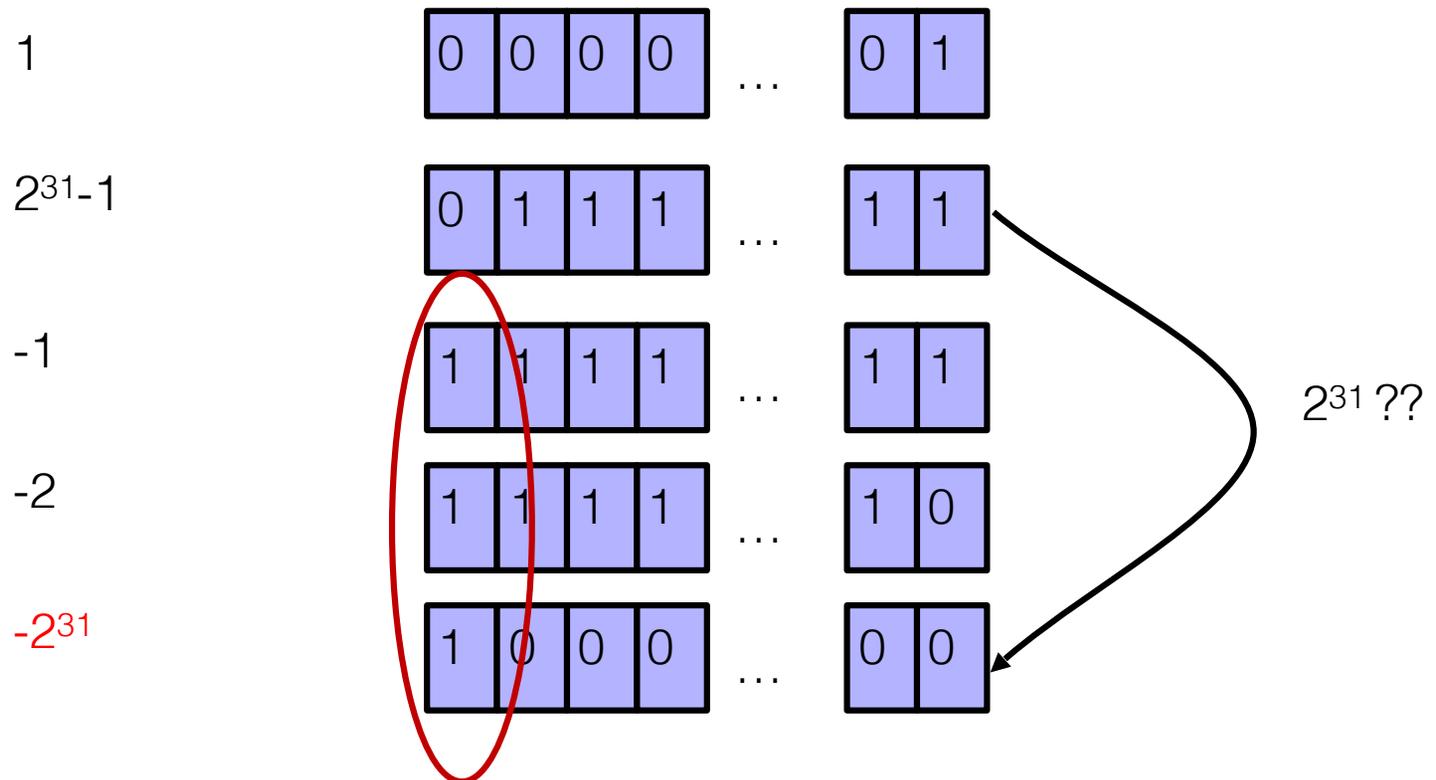
Integer Overflow Attacks

Two's Complement

Binary representation of negative integers

Represent X (where $X < 0$) as $2^N - |X|$

N is word size (e.g., 32 bits on x86 architecture)



Integer Overflow

```
static int getpeername1(p, uap, compat) {
```

```
// In FreeBSD kernel, retrieves address of peer to which a socket is connected
```

```
...
```

```
struct sockaddr *sa;
```

Checks that "len" is not too big

```
...
```

Negative "len" will always pass this check...

```
assert(len = min(len, sa->sa_len));
```

```
copyout(sa, (caddr_t)uap->asa, (u_int)len);
```

```
...
```

... interpreted as a huge unsigned integer here

```
}
```

Copies "len" bytes from kernel memory to user space

... will copy up to 4G of kernel memory

Format String Attacks

Variable Arguments in C

- ◆ In C, can define a function with a variable number of arguments
 - Example: `void printf(const char* format, ...)`
- ◆ Examples of usage:

```
printf("hello, world");  
printf("length of %s = %d\n", str, str.length());  
printf("unable to open file descriptor %d\n", fd);
```

Format specification encoded by special % characters

%d,%i,%o,%u,%x,%X – integer argument
%s – string argument
%p – pointer argument (void *)
Several others

Implementation of Variable Args

Special functions `va_start`, `va_arg`, `va_end` compute arguments at run-time

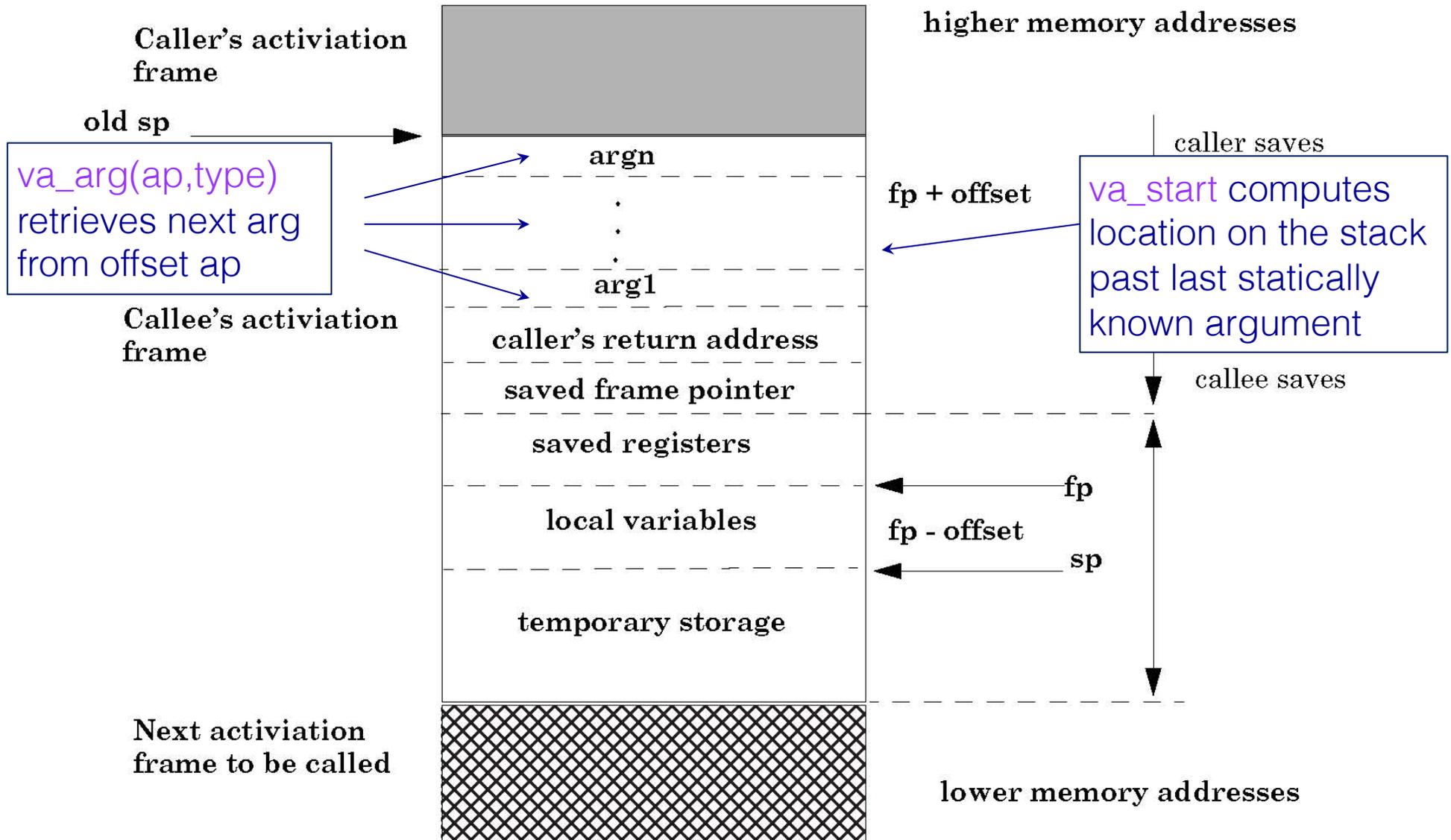
```
void printf(const char* format, ...)
{
    int i; char c; char* s; double d;
    va_list ap; /* declare an "argument pointer" to a variable arg list */
    va_start(ap, format); /* initialize arg pointer using last known arg */

    for (char* p = format; *p != '\\0'; p++) {
        if (*p == '%') {
            switch (*++p) {
                case 'd':
                    i = va_arg(ap, int); break;
                case 's':
                    s = va_arg(ap, char*); break;
                case 'c':
                    c = va_arg(ap, char); break;
            }
            ... /* etc. for each % specification */
        }
    }
    ...

    va_end(ap); /* restore any special stack manipulations */
}
```

printf has an internal stack pointer

Frame with Variable Args



Format Strings in C

◆ Proper use of printf format string:

```
... int foo=1234;  
    printf("foo = %d in decimal, %X in hex",foo,foo); ...
```

This will print

```
foo = 1234 in decimal, 4D2 in hex
```

◆ Sloppy use of printf format string:

```
... char buf[13]="Hello, world!";  
    printf(buf);  
    // should've used printf("%s", buf); ...
```

If the buffer contains a format symbol starting with %, location pointed to by printf's internal stack pointer will be interpreted as an argument of printf. This can be exploited to **move printf's internal stack pointer!** (how?)

Writing Stack with Format Strings

- ◆ `%n` format symbol tells `printf` to write the number of characters that have been printed

```
... printf("Overflow this!%n", &myVar); ...
```

Argument of `printf` is interpreted as destination address

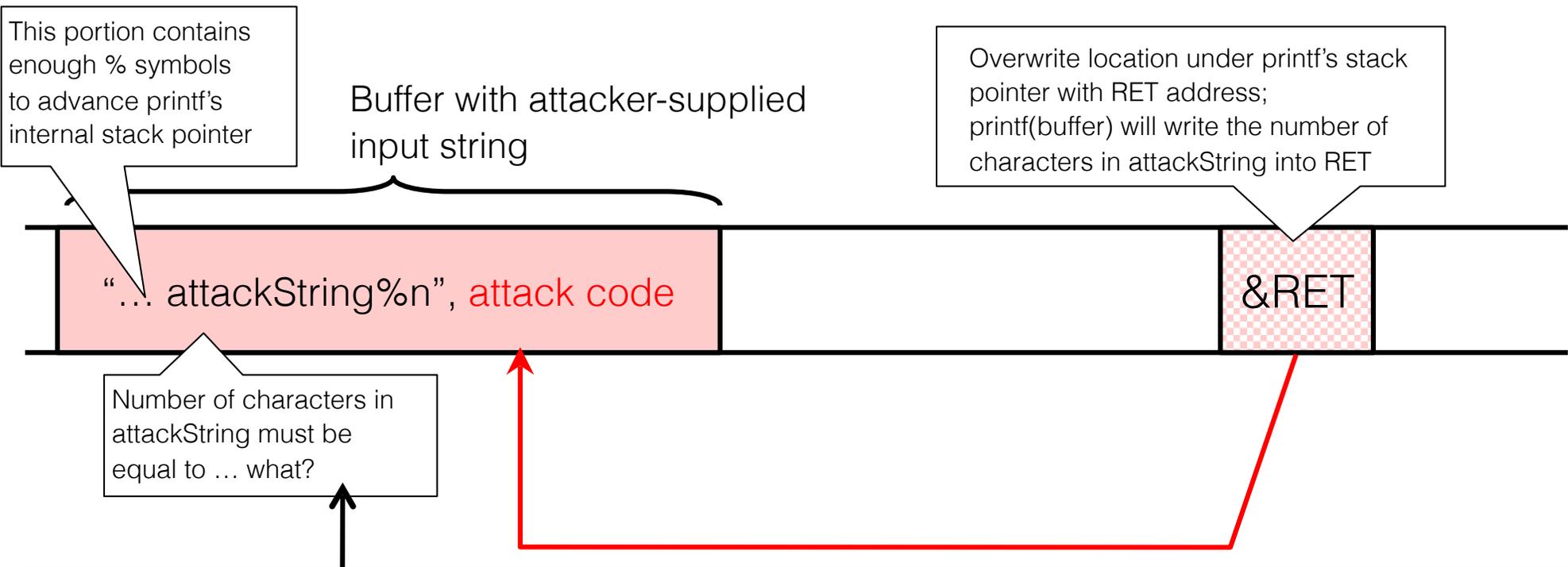
This writes `14` into `myVar` ("Overflow this!" has 14 characters)

- ◆ What if `printf` does not have an argument?

```
... char buf[16]="Overflow this!%n";  
printf(buf); ...
```

Stack location pointed to by `printf`'s internal stack pointer will be interpreted as address into which the number of characters will be written!

Using %n to Mung Return Address



C has a concise way of printing multiple symbols: %Mx will print exactly 4M bytes (taking them from the stack). Attack string should contain enough "%Mx" so that the number of characters printed is equal to the most significant byte of the address of the attack code.

See "Exploiting Format String Vulnerabilities" for details